

Technical Report

Analysis of the Convergence Rate of SoPro

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In this report, we discuss the impact of the optimization problem and the network topology on the convergence rate of SoPro [1]. Please see [1] for all the assumptions and notations.

Theorem 1 in [1] says that the linear convergence rate of SoPro is given by

$$\|\mathbf{z}^{k+1} - \mathbf{z}^*\|_Q^2 \leq (1 - \delta) \|\mathbf{z}^k - \mathbf{z}^*\|_Q^2,$$

where

$$\delta = \sup_{c_1, c_2 > 0} \min \left\{ \frac{\rho \lambda_W \kappa_{\beta, \eta}}{(1 + c_1) \|\Lambda_M + D\|^2}, \frac{1 - \eta}{(1 + 1/c_1)(1 + c_2)}, \frac{2\eta m_\rho - \beta}{\lambda_{\max}(R + (1 + 1/c_1)(1 + 1/c_2)\Lambda_M^2 / (\rho \lambda_W))} \right\}. \quad (1)$$

Here, λ_W is the smallest nonzero eigenvalue of the weight matrix W and $\beta \in (0, 2\eta m_\rho)$ is such that $\kappa_{\beta, \eta} := \lambda_{\min}(R - \frac{\Lambda_M}{2(1-\eta)} - \frac{(\Lambda_M - \Lambda_m)^2}{4\beta} - \Lambda_M + \Lambda_m - \rho W) > 0$.

Below, we analyze the convergence rate under the following setting:

- 1) Each f_i , $i \in \mathcal{V}$ is globally strongly convex with convexity parameter $m_i > 0$.
- 2) Without loss of generality, we set $m_i = m \forall i \in \mathcal{V}$ and $M_i = M \forall i \in \mathcal{V}$ for some $0 < m < M$. Also, let $m_\rho = m$.
- 3) For simplicity, we choose $W = I_{Nd} - A \otimes I_d$, where $A \in \mathbb{R}^{n \times n}$ is a symmetric doubly stochastic matrix so that $\|A\| \leq 1$ and thus $\|W\| \leq 2$. To satisfy the condition on W in [1], we require $A_{ij} > 0$ if $i = j$ or $\{i, j\} \in \mathcal{E}$ and $A_{ij} = 0$ otherwise.
- 4) Let $\eta = \frac{1}{2}$ and arbitrarily pick $\rho, \epsilon > 0$.
- 5) Let $D = \rho W + (\frac{M}{2(1-\eta)} + \frac{(M-m)^2}{4\eta m} + \frac{M-3m}{2} + \epsilon)I_{Nd} = \rho W + (\frac{(M-m)^2}{2m} + \frac{3(M-m)}{2} + \epsilon)I_{Nd}$, which satisfies the condition on D in Lemma 2.
- 6) Let $\beta = m_\rho \eta = \frac{m}{2}$ so that $\kappa_{\beta, \eta} = \epsilon > 0$.

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With the above setting, note that $\lambda_{\max}(D) \leq \frac{(M-m)^2}{2m} + \frac{3(M-m)}{2} + \epsilon + 2\rho =: c$. Thus, for any $c_1, c_2 > 0$,

$$\begin{aligned} \frac{\rho\lambda_W\kappa_{\beta,\eta}}{(1+c_1)\|\Lambda_M + D\|^2} &\geq \frac{\rho\epsilon\lambda_W}{(1+c_1)(M+c)^2} =: Q_1(c_1, c_2), \\ \frac{1-\eta}{(1+1/c_1)(1+c_2)} &= \frac{1}{2(1+1/c_1)(1+c_2)} =: Q_2(c_1, c_2), \\ \frac{2\eta m_\rho - \beta}{\lambda_{\max}(R + (1+1/c_1)(1+1/c_2)\Lambda_M^2/(\rho\lambda_W))} &\geq \frac{m}{2c + M + m + 2(1+1/c_1)(1+1/c_2)M^2/(\rho\lambda_W)} \\ &:= Q_3(c_1, c_2). \end{aligned}$$

Therefore,

$$\|\mathbf{z}^{k+1} - \mathbf{z}^*\|_Q^2 \leq (1 - \hat{\delta})\|\mathbf{z}^k - \mathbf{z}^*\|_Q^2,$$

where

$$\hat{\delta} = \sup_{c_1, c_2 > 0} \min\{Q_1(c_1, c_2), Q_2(c_1, c_2), Q_3(c_1, c_2)\} \leq \delta.$$

Note that for any $c_1, c_2 > 0$, $\min\{Q_1(c_1, c_2), Q_2(c_1, c_2), Q_3(c_1, c_2)\} \in (0, 1)$. Hence, $0 < \hat{\delta} < 1$ and we can use $\hat{\delta}$ to bound the convergence rate of SoPro.

It can be seen that $\hat{\delta}$ only depends on m , M , and λ_W . Below, we discuss how these factors affect the value of $\hat{\delta}$. To this end, consider the following lemma:

Lemma 1. *Suppose $\sup_{c_1, c_2 > 0} \min\{Q_1(c_1, c_2), Q_2(c_1, c_2), Q_3(c_1, c_2)\} = \min\{Q_1(c_1^*, c_2^*), Q_2(c_1^*, c_2^*), Q_3(c_1^*, c_2^*)\}$ for some $c_1^*, c_2^* > 0$. Then, $Q_1(c_1^*, c_2^*) = Q_2(c_1^*, c_2^*) = Q_3(c_1^*, c_2^*)$.*

Proof. See Appendix A. □

Based on Lemma 1, we derive the lemma below, which provides a sufficient condition on the strict increase of $\hat{\delta}$.

Lemma 2. *Suppose $Q_i(c_1, c_2) \forall i = 1, 2, 3$ are associated with problem \mathcal{P} and network \mathcal{G} . Let $Q'_i(c_1, c_2) \forall i = 1, 2, 3$ be defined in the same way as $Q_i(c_1, c_2) \forall i = 1, 2, 3$ but correspond to a different problem \mathcal{P}' and a different network \mathcal{G}' . If (i) $Q'_i(c_1, c_2) \geq Q_i(c_1, c_2) \forall i = 1, 2, 3$ and (ii) there exists $j \in \{1, 2, 3\}$ such that $Q'_j(c_1, c_2) > Q_j(c_1, c_2)$, then*

$$\sup_{c_1, c_2 > 0} \min\{Q'_1(c_1, c_2), Q'_2(c_1, c_2), Q'_3(c_1, c_2)\} > \sup_{c_1, c_2 > 0} \min\{Q_1(c_1, c_2), Q_2(c_1, c_2), Q_3(c_1, c_2)\}. \quad (2)$$

Proof. See Appendix B. □

For any $c_1, c_2 > 0$, observe that

- 1) $Q_1(c_1, c_2)$ and $Q_3(c_1, c_2)$ increase if m increases, M decreases, or λ_W increases.

2) $Q_2(c_1, c_2)$ is independent of m , M , and λ_W .

From Lemma 2, we thereby conclude that **larger m , smaller M , or larger λ_W leads to larger $\hat{\delta}$, which implies faster convergence of SoPro.** Here, m and M are the problem characteristics and λ_W indicates how well the network is connected (which plays a similar role as the algebraic connectivity).

APPENDIX

A. Proof of Lemma 1

Equivalently, we show that given $c_1, c_2 > 0$, if $Q_i(c_1, c_2) \forall i = 1, 2, 3$ are nonidentical, then there exist $c'_1, c'_2 > 0$ such that

$$\min\{Q_1(c'_1, c'_2), Q_2(c'_1, c'_2), Q_3(c'_1, c'_2)\} > \min\{Q_1(c_1, c_2), Q_2(c_1, c_2), Q_3(c_1, c_2)\}. \quad (3)$$

To this end, *suppose $Q_i(c_1, c_2) \forall i = 1, 2, 3$ are nonidentical* and consider the following three mutually exclusive and exhaustive cases:

Case 1: $Q_1(c_1, c_2) = \min\{Q_2(c_1, c_2), Q_3(c_1, c_2)\}$, where $Q_2(c_1, c_2) \neq Q_3(c_1, c_2)$.

We execute the following two steps to find $c'_1, c'_2 > 0$ such that (3) holds:

Step 1: Find $c'_2 > 0$ such that $\min\{Q_2(c_1, c'_2), Q_3(c_1, c'_2)\} > \min\{Q_2(c_1, c_2), Q_3(c_1, c_2)\}$. To do so, consider two subcases as follows:

- If $Q_2(c_1, c_2) < Q_3(c_1, c_2)$, then decrease c_2 to some $c'_2 \in (0, c_2)$ satisfying $Q_2(c_1, c'_2) = Q_3(c_1, c'_2)$.
- If $Q_2(c_1, c_2) > Q_3(c_1, c_2)$, then increase c_2 to some $c'_2 \in (c_2, \infty)$ satisfying $Q_2(c_1, c'_2) = Q_3(c_1, c'_2)$.

Because $Q_2(c_1, c_2)$ increases with the decrease of c_2 and goes to 0 as $c_2 \rightarrow \infty$ and because $Q_3(c_1, c_2)$ increases with the increase of c_2 and goes to 0 as $c_2 \rightarrow 0$, c'_2 in Step 1 exists. Also, since Q_1 is independent of c_2 , $Q_1(c_1, c'_2) = Q_1(c_1, c_2) = \min\{Q_2(c_1, c_2), Q_3(c_1, c_2)\} < \min\{Q_2(c_1, c'_2), Q_3(c_1, c'_2)\}$.

Step 2: Find $c'_1 > 0$ such that $Q_1(c_1, c'_2) < Q_1(c'_1, c'_2) = \min\{Q_2(c'_1, c'_2), Q_3(c'_1, c'_2)\}$. To do so, go to Case 2 (with c_2 replaced by c'_2).

With the above two steps, $\min\{Q_1(c'_1, c'_2), Q_2(c'_1, c'_2), Q_3(c'_1, c'_2)\} = Q_1(c'_1, c'_2) > Q_1(c_1, c'_2) = Q_1(c_1, c_2) = \min\{Q_1(c_1, c_2), Q_2(c_1, c_2), Q_3(c_1, c_2)\}$, i.e., (3) holds.

Case 2: $Q_1(c_1, c_2) < \min\{Q_2(c_1, c_2), Q_3(c_1, c_2)\}$. To meet (3), decrease c_1 to some $c'_1 \in (0, c_1)$ satisfying $Q_1(c'_1, c_2) = \min\{Q_2(c'_1, c_2), Q_3(c'_1, c_2)\}$. Because $Q_1(c_1, c_2)$ increases with the decrease of c_1 and goes to 0 as $c_1 \rightarrow \infty$ and because $Q_2(c_1, c_2)$ and $Q_3(c_1, c_2)$ decrease with the decrease of c_1 and both go to 0 as $c_1 \rightarrow 0$, such c'_1 exists. Since $\min\{Q_1(c_1, c_2), Q_2(c_1, c_2), Q_3(c_1, c_2)\} = Q_1(c_1, c_2) < Q_1(c'_1, c_2) = \min\{Q_1(c'_1, c_2), Q_2(c'_1, c_2), Q_3(c'_1, c_2)\}$, (3) holds by letting $c'_2 = c_2$.

Case 3: $Q_1(c_1, c_2) > \min\{Q_2(c_1, c_2), Q_3(c_1, c_2)\}$. In this case, increase c_1 to some $c'_1 \in (c_1, \infty)$ satisfying $Q_1(c'_1, c_2) = \min\{Q_2(c'_1, c_2), Q_3(c'_1, c_2)\}$. Similar to Case 2, such c'_1 exists. Since

$\min\{Q_1(c'_1, c_2), Q_2(c'_1, c_2), Q_3(c'_1, c_2)\} = \min\{Q_2(c'_1, c_2), Q_3(c'_1, c_2)\} > \min\{Q_2(c_1, c_2), Q_3(c_1, c_2)\} = \min\{Q_1(c_1, c_2), Q_2(c_1, c_2), Q_3(c_1, c_2)\}$, (3) holds by letting $c'_2 = c_2$.

From the above three cases, for any $c_1, c_2 > 0$, as long as $Q_i(c_1, c_2) \forall i = 1, 2, 3$ are not identical, the value of $\min\{Q_1(c_1, c_2), Q_2(c_1, c_2), Q_3(c_1, c_2)\}$ can be increased by tuning c_1 and c_2 . This completes the proof.

B. Proof of Lemma 2

Let $c_1^*, c_2^* > 0$ be such that

$$\min\{Q_1(c_1^*, c_2^*), Q_2(c_1^*, c_2^*), Q_3(c_1^*, c_2^*)\} = \sup_{c_1, c_2 > 0} \min\{Q_1(c_1, c_2), Q_2(c_1, c_2), Q_3(c_1, c_2)\}.$$

From Lemma 1, $Q_1(c_1^*, c_2^*) = Q_2(c_1^*, c_2^*) = Q_3(c_1^*, c_2^*)$. Suppose conditions (i) and (ii) in the statement of Lemma 2 hold. Consider the following two mutually exclusive and exhaustive cases:

- $Q'_i(c_1^*, c_2^*) > Q_i(c_1^*, c_2^*) \forall i = 1, 2, 3$. In this case, $\sup_{c_1, c_2 > 0} \min\{Q'_1(c_1, c_2), Q'_2(c_1, c_2), Q'_3(c_1, c_2)\} \geq \min\{Q'_1(c_1^*, c_2^*), Q'_2(c_1^*, c_2^*), Q'_3(c_1^*, c_2^*)\} > \min\{Q_1(c_1^*, c_2^*), Q_2(c_1^*, c_2^*), Q_3(c_1^*, c_2^*)\}$.
- There exists $i \in \{1, 2, 3\}$ such that $Q'_i(c_1^*, c_2^*) = Q_i(c_1^*, c_2^*)$. In this case, due to condition (ii), there exists $j \in \{1, 2, 3\} \setminus \{i\}$ such that $Q'_j(c_1^*, c_2^*) > Q_j(c_1^*, c_2^*) = Q_i(c_1^*, c_2^*)$. Thus, $Q'_j(c_1^*, c_2^*) \neq Q'_i(c_1^*, c_2^*)$. It follows from Lemma 1 that

$$\sup_{c_1, c_2 > 0} \min\{Q'_1(c_1, c_2), Q'_2(c_1, c_2), Q'_3(c_1, c_2)\} > \min\{Q'_1(c_1^*, c_2^*), Q'_2(c_1^*, c_2^*), Q'_3(c_1^*, c_2^*)\}.$$

In both of the above cases, (2) holds.

REFERENCES

- [1] X. Wu, Z. Qu, and J. Lu, "A second-order proximal algorithm for consensus optimization," 2019, submitted to *IEEE Transactions on Automatic Control*.