



# Discussion 4 RISC-V

# Agenda



上海科技大学  
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1. Big Endian vs. Little Endian
2. Label and Assembler Directives
3. Enviroment calls
4. RISC-V Practices
5. Venus Q&A



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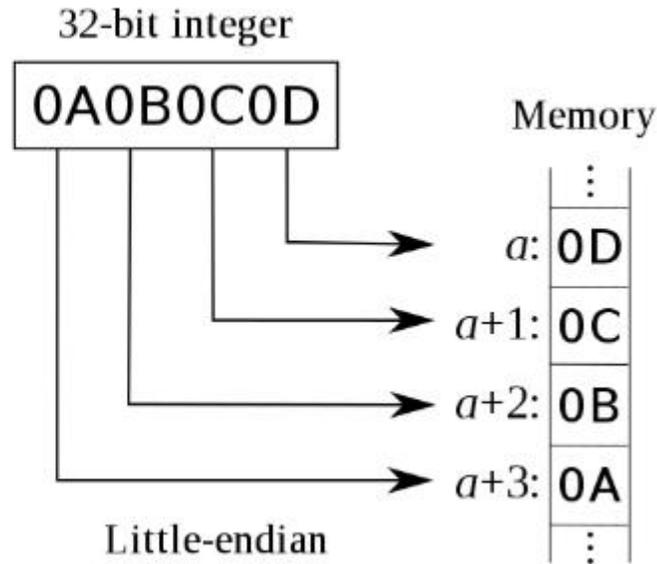
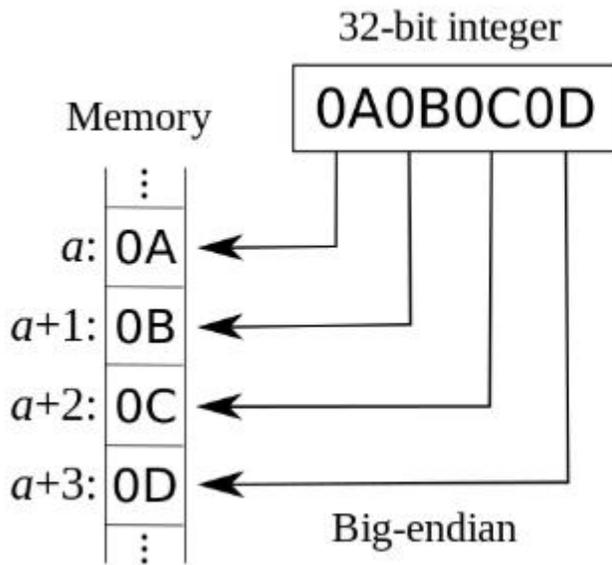
# Big Endian vs. Little Endian

MSB : the highest-order bit in a binary number. It represents the most significant part of the value.

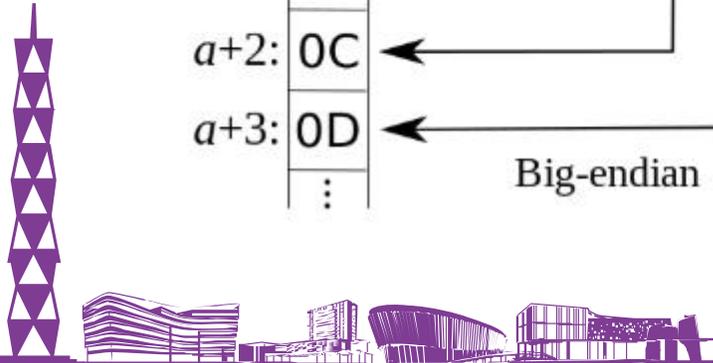
LSB : the lowest-order bit in a binary number. It represents the least significant part of the value.

Big-Endian: MSB at lowest address

Little-Endian: LSB at lowest address



Example:  
 0x0A0B0C0D  
 MSB = 0x0A  
 LSB = 0x0D





Directives are instructions to the assembler that tell it how to organize code, data, and symbols.

Label: a label is a symbolic name for the address of a data item or an instruction.

- `.text`: Subsequent items put in user text segment (machine code)
- `.data`: Subsequent items put in user data segment (binary rep of data in source file)
- `.globl sym`: declares `sym` global and can be referenced from other files
- `.ascii str`: Store the string `str` in memory and null-terminates it
- `.word w1...wn`: Store the  $n$  32-bit quantities in successive memory words

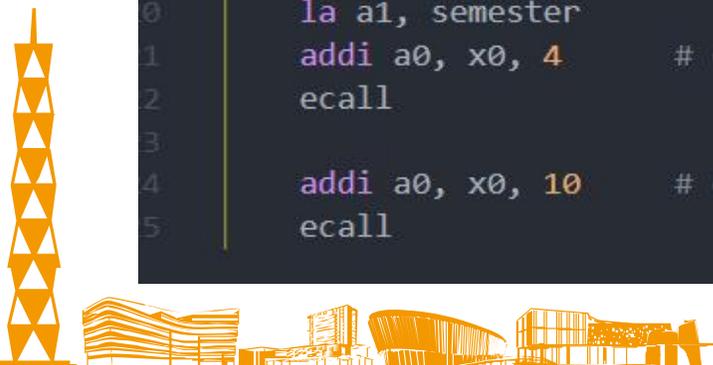


```
1 .data
2 course:
3   .ascii "cs110"
4
5 semester:
6   .ascii "sp26"
7
8 num:
9   .word 2026
10
11 .text
12 la a1, course
13 addi a0, x0, 4 # ecall 4 -- print_string
14 ecall
15
16 addi a1, x0, 10 # ASCII 10 -- '\n'
17 addi a0, x0, 11 # ecall 11 -- print_character
18 ecall
19
20 la a1, semester
21 addi a0, x0, 4 # ecall 4 -- print_string
22 ecall
23
24 addi a0, x0, 10 # ecall 10 -- exit
25 ecall
```

label : course , semester,num

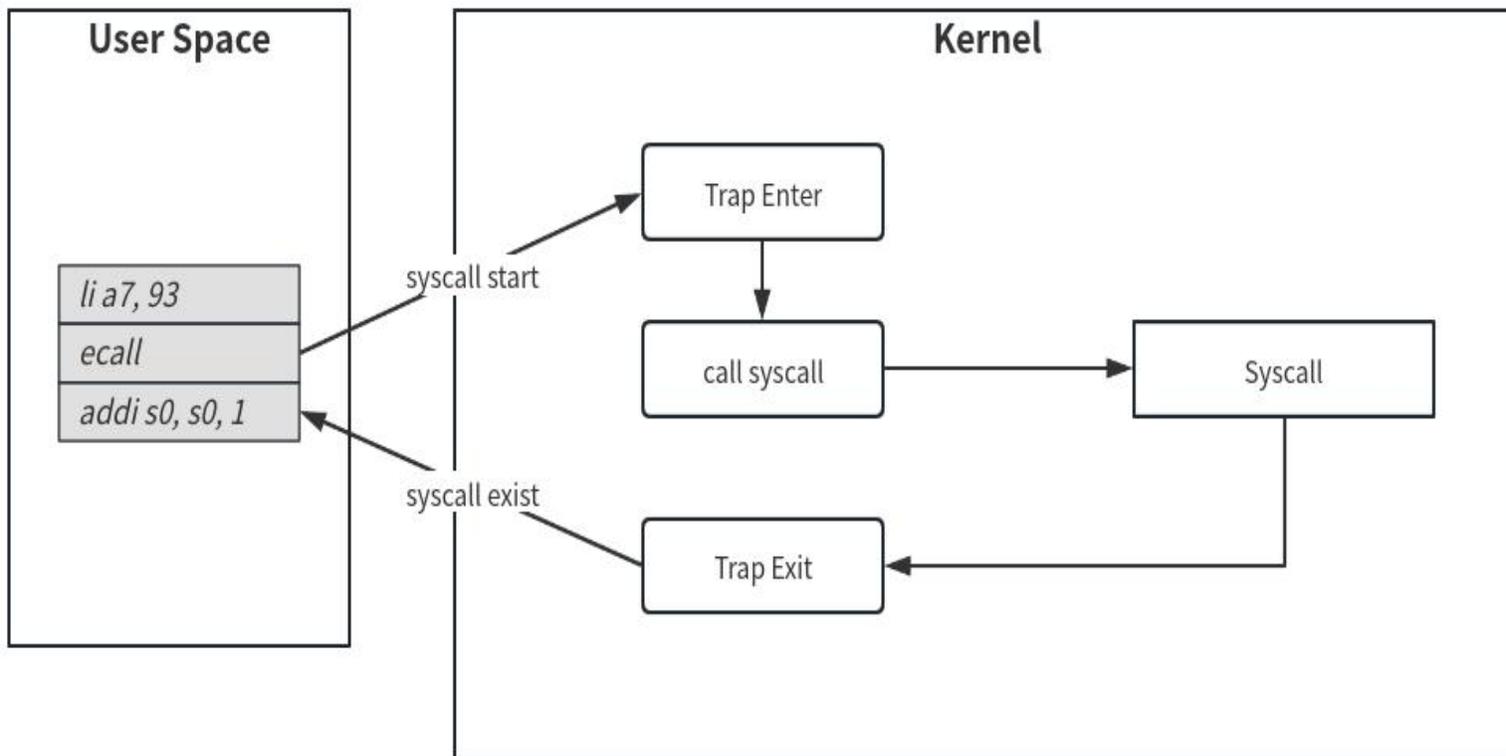
assembler directives

- section management : data,text
- data definition : ascii , word





# Ecall



Dual mode

- User program runs in user model
- OS services run in kernel mode

Triggers a trap

- Switches to kernel mode
- Kernel handles the request
- Returns to user program





```
1  .data
2  course:
3      .asciiiz "cs110"
4
5  semester:
6      .asciiiz "sp26"
7
8  num:
9      .word 2026
10
11 .text
12  la a1, course
13  addi a0, x0, 4      # ecall 4 -- print_string
14  ecall
15
16  addi a1, x0, 10     # ASCII 10 -- '\n'
17  addi a0, x0, 11     # ecall 11 -- print_character
18  ecall
19
20  la a1, semester
21  addi a0, x0, 4      # ecall 4 -- print_string
22  ecall
23
24  addi a0, x0, 10     # ecall 10 -- exit
25  ecall
```

## How to Use ecall

- a0 → syscall number
- a1 → argument

## Output

```
cs110
sp26|
```



## 4. RISC-V

- 8 (a) Consider the following code snippet written in RISC-V. The function `Factorial` is to calculate the factorial of a given number. (i.e.  $n! = n \cdot (n - 1) \cdots 2 \cdot 1$ )

```
1 Factorial:
2   addi   sp, sp, -8
3   sw     ra, 0(sp)
4   li     t0, 1
5   beq    a0, t0, last_sit
6   sw     a0, 4(sp)
7   _____
8   _____
9   lw     t0, 4(sp)
10  _____
11  j      fact_done
12 last_sit:
13  _____
14 fact_done:
15  lw     ra, 0(sp)
16  addi   sp, sp, 8
17  mv     a0, a1
18  jr     ra
```

Fill in the missing code below.

line 7: \_\_\_\_\_

line 8: `jal Factorial` \_\_\_\_\_

line 10: \_\_\_\_\_

line 13: \_\_\_\_\_

## 4. RISC-V

- 8 (a) Consider the following code snippet written in RISC-V. The function `Factorial` is to calculate the factorial of a given number. (i.e.  $n! = n \cdot (n - 1) \cdots 2 \cdot 1$ )

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2   addi   sp, sp, -8
3   sw     ra, 0(sp)
4   li     t0, 1
5   beq    a0, t0, last_sit
6   sw     a0, 4(sp)
7   _____
8   _____
9   lw     t0, 4(sp)
10  _____
11  j      fact_done
12 last_sit:
13  _____
14 fact_done:
15  lw     ra, 0(sp)
16  addi   sp, sp, 8
17  mv     a0, a1
18  jr     ra
```

Fill in the missing code below.

line 7: \_\_\_\_\_

line 8: \_\_\_\_\_

line 10: \_\_\_\_\_

line 13: \_\_\_\_\_

### Solution:

line 7: `addi a0, a0, -1,`

line 8: `jal Factorial,`

line 10: `mul a1, t0, a1,`

line 13: `li a1, 1` or `addi a1, x0, 1.`

# RISC-V Practices



**Singly-linked list** is a common and useful data structure. In this problem, you are going to implement a linked list operation in RISC-V assembly. Assume the assembly is for a **32-bit machine**. Also, by convention, consecutive fields occupy consecutive bytes within the structure by their declaration order, and the first field takes the lowest address. The node in a singly-linked list is defined as a [struct](#) type as follows.

```
struct Node
{
    // Value of this node
    int val;
    // Pointer to the next node
    struct Node *next_node;
};
```

```
1 # a0: address of node A; a1: address of node B
2 insert_node:
3     lw t0 4(a0)
4
5
6     ret
```

Then we define the function:

- insert node : Given a pointer to node A and a pointer to node B, this function will insert node B into the linked list, making node B the next node of node A. Node A is already in the list and assume that it is not the last node (tail) of the list.

**Singly-linked list** is a common and useful data structure. In this problem, you are going to implement two linked list operations in RISC-V assembly. Assume the assembly is for a **32-bit machine**. Also, by convention, consecutive fields occupy consecutive bytes within the structure by their declaration order, and the first field takes the lowest address. The node in a singly-linked list is defined as a [struct](#) type as follows.

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- insert node : Given a pointer to node A and a pointer to node B, this function will insert node B into the linked list, making node B the next node of node A. Node A is already in the list and assume that it is not the last node (tail) of the list.

```
1 # a0: address of node A; a1: address of node B
2 insert_node:
3     # temp = A->next
4     lw t0 4(a0)
5     # B->next = temp
6     sw t0 4(a1)
7     # A->next = B
8     sw a1 4(a0)
9     ret
```

# RISC-V Practices



(a) Doubly linked list is a common and useful data structure. In this problem, you are going to implement two linked list operations in RISC-V assembly. Assume the assembly is for a 32-bit machine. Also, by convention, consecutive fields occupy consecutive bytes within the structure by their declaration order, and the first field takes the lowest address. The node in a double linked list is defined as a `struct` type as follows.

```
1 struct node{
2     // value of this node
3     int val;
4     // pointer to next node
5     struct node * next_node;
6     // pointer to previous node
7     struct node * prev_node;
8 };
```

Then we define some functions:

- `insert_node` : Given a pointer to node A and a pointer to node B, this function will insert node B into the linked list, making node B the next node of node A. Node A is already in the list and assume that it is not the last node (tail) of the list.
- `switch_node` : Given a pointer to node A and a pointer to node B (A and B are different and they are not adjacent), this function will exchange the location of node A and node B in the linked list without changing the node values. Assume that nodes A and B are neither the head nor the tail of the linked list, otherwise, they can be at any positions in the linked list.

Please fill in the following RISC-V codes to implement these two functions

```
// a0: address of node A; a1: address of node B
insert_node:
    lw t0 4(a0)
```

```
_____
_____
_____
_____
ret
```

```
// a0: address of node A; a1: address of node B
switch_node:
    lw t0 4(a0)
    lw t1 4(a1)
    lw t2 8(a0)
    lw t3 8(a1)
```

```
_____
_____
_____
_____
_____
_____
ret
```

- (a) Doubly linked list is a common and useful data structure. In this problem, you are going to implement two linked list operations in RISC-V assembly. Assume the assembly is for a 32-bit machine. Also, by convention, consecutive fields occupy consecutive bytes within the structure by their declaration order, and the first field takes the lowest address. The node in a double linked list is defined as a `struct` type as follows.

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6     // pointer to previous node
7     struct node * prev_node;
8 };
```

Then we define some functions:

- `insert_node` : Given a pointer to node A and a pointer to node B, this function will insert node B into the linked list, making node B the next node of node A. Node A is already in the list and assume that it is not the last node (tail) of the list.
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Please fill in the following RISC-V codes to implement these two functions

```
// a0: address of node A; a1: address of
// node B
insert_node
// t0 for A->next_node
lw t0 4(a0)
// B->next_node = A->next_node
sw t0 4(a1)
// A->next_node = B
sw a1 4(a0)
// B->prev_node = A
sw a0 8(a1)
// B->next_node->prev_node = B
sw a1 8(t0)
ret
```

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- `insert_node` : Given a pointer to node A and a pointer to node B, this function will insert node B into the linked list, making node B the next node of node A. Node A is already in the list and assume that it is not the last node (tail) of the list.
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Please fill in the following RISC-V codes to implement these two functions

```
switch_node:
// temp1 = A->next
// temp2 = B->next
// temp3 = A->prev
// temp4 = B->prev
// A->next->prev = B
// A->prev->next = B
// B->next->prev = A
// B->prev->next = A
// B->prev = A->prev
// B->next = A->next
// A->prev = B->prev
// A->next = B->next
```

```
lw t0 4(a0)
lw t1 4(a1)
lw t2 8(a0)
lw t3 8(a1)
sw a1 8(t0)
sw a1 4(t2)
sw a0 8(t1)
sw a0 4(t3)
sw t2 8(a1)
sw t0 4(a1)
sw t3 8(a0)
sw t1 4(a0)
```

# RISC-V Practices



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```
1 # a0: address of node A; a1: address of node B
2 insert_node:
3     lw t1 4(a1)
4     lw t2 8(a0)
5
6     ret
```

Consider another scenario for switch node:

- switch node : Given a pointer to node A and a pointer to node B (**A and B are different and they are adjacent, the next node of A is B**), this function will exchange the location of node A and node B in the linked list without changing the node values. Assume that nodes A and B are neither the head nor the tail of the linked list, otherwise, they can be at any positions in the linked list.

# RISC-V Practices



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```
1 # a0: address of node A; a1: address of node B
2 insert_node:
3     # temp1 = B->next
4     lw t1 4(a1)
5     # temp2 = A->prev
6     lw t2 8(a0)
7     # B->next->prev = A
8     sw a0 8(t1)
9     # A->prev->next = B
10    sw a1 4(t2)
11    # A->next = B->next
12    sw t1 4(a0)
13    # A->prev = B
14    sw a1 8(a0)
15    # B->next = A
16    sw a0 4(a1)
17    # B->prev = A->prev
18    sw t2 8(a1)
19    ret
```



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# Venus Q&A



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Thank you for attending the discussion!

Wish you good luck doing homework/projects/exams!



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